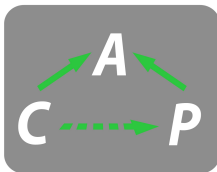


How to compute in free abelian categories

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Theorem

Let A be a commutative \mathbb{Q} -algebra, $a, b \in A$ such that

$$a^3b + a = b^3a + b = a^2b^2 + 1 = 0.$$

Then $1 = 0$ in A .

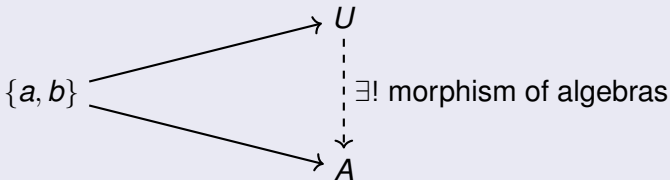
Proof strategy

- 1 Construct the universal commutative algebra U which satisfies the premises of the theorem.
- 2 Show that the conclusion $1 = 0$ holds for U .

- 1 Construct the universal commutative algebra U which satisfies the premises of the theorem.

The universal algebra which satisfies the premises.

- a commutative algebra U
- a map $\{a, b\} \rightarrow U$ with $a^3b + a = b^3a + b = a^2b^2 + 1 = 0$.
- For every commutative algebra A and map $\{a, b\} \rightarrow A$ whose images in A also satisfy that relation:



$$U := \mathbb{Q}[a, b] / \langle a^3b + a, b^3a + b, a^2b^2 + 1 \rangle$$

- 2 Show that the conclusion $1 = 0$ holds for U . (Gröbner basis)

Q: Can we apply a similiar proof strategy for the Snake lemma?

Snake lemma

Suppose given the following commutative diagram with exact rows in an abelian category:

$$\begin{array}{ccccccc} & & A & \xrightarrow{a} & B & \longrightarrow & \text{Cok}(a) & \longrightarrow & 0 \\ & & \downarrow e & & \downarrow b & & \downarrow d & & \\ 0 & \longrightarrow & \text{Ker}(c) & \longrightarrow & C & \xrightarrow{c} & D & & \end{array}$$

Then we have an exact sequence

$$\text{Ker}(e) \longrightarrow \text{Ker}(b) \longrightarrow \text{Ker}(d) \xrightarrow{\partial} \text{Cok}(e) \longrightarrow \text{Cok}(b) \longrightarrow \text{Cok}(d)$$

1 Free abelian categories

2 A computational proof of the Snake lemma

1 Free abelian categories

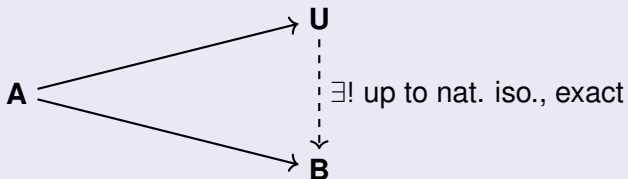
2 A computational proof of the Snake lemma

Free abelian categories

Let \mathbf{A} be an additive category.

Data of the free abelian category

- an abelian category \mathbf{U}
- an additive functor $\mathbf{A} \rightarrow \mathbf{U}$
- For every additive functor $\mathbf{A} \rightarrow \mathbf{B}$ into an abelian category:



How to construct free abelian categories

- Existence: Peter Freyd
- Explicit construction: Murray Adelman

Homology of sequences

How to compute homology in an abelian category:

$$\begin{array}{ccccc} A & \xrightarrow{a} & B & \xrightarrow{b} & C \\ & & \nearrow & \nwarrow & \\ & & \text{Im}(a) & \xrightarrow{e} & \text{Ker}(b) \twoheadrightarrow \text{Cok}(e) \end{array}$$

This construction only makes sense if we have $a \cdot b = 0$.

Homology of composite morphisms

How to compute homology in an abelian category:

$$\begin{array}{ccccc} A & \xrightarrow{a} & B & \xrightarrow{b} & C \\ & & \nearrow & & \searrow \\ & & \text{Ker}(b) & \xrightarrow{\quad} & \text{Cok}(a) \\ & & \searrow & & \nearrow \\ & & & & H(A \xrightarrow{a} B \xrightarrow{b} C) \end{array}$$

This construction makes sense even if we don't have $a \cdot b = 0$.

Adelman's observation

Equipping an additive category with "homologies" makes it abelian.

Some examples of morphisms

\mathbf{A} = Category of free \mathbb{Z} -modules

$$\begin{array}{ccccc} (\mathbb{Z} & \xrightarrow{4} & \mathbb{Z} & \longrightarrow & 0) \\ \downarrow 2 & \swarrow s_1 & \downarrow 1 & \searrow s_2 & \downarrow \\ (\mathbb{Z} & \xrightarrow{2} & \mathbb{Z} & \longrightarrow & 0) \end{array}$$

A non-zero morphism in $\text{Adel}(\mathbf{A})$.

This reminds us of a presentation of the morphism $\mathbb{Z}/4\mathbb{Z} \rightarrow \mathbb{Z}/2\mathbb{Z}$ by free modules.

Some examples of morphisms

\mathbf{A} = Category of free \mathbb{Z} -modules

$$\begin{array}{ccccc} (0 & \longrightarrow & \mathbb{Z} & \xrightarrow{2} & \mathbb{Z}) \\ \downarrow \text{---} & & \downarrow 1 & & \downarrow 2 \text{---} \\ (0 & \longrightarrow & \mathbb{Z} & \xrightarrow{4} & \mathbb{Z}) \end{array}$$

A non-zero morphism in $\text{Adel}(\mathbf{A})$.

Some examples of morphisms

\mathbf{A} = Category of free \mathbb{Z} -modules

$$\begin{array}{ccccc} (\mathbb{Z} & \xrightarrow{2} & \mathbb{Z} & \xrightarrow{2} & \mathbb{Z}) \\ \downarrow & & \downarrow 1 & & \downarrow \\ (\mathbb{Z} & \xrightarrow{2} & \mathbb{Z} & \xrightarrow{2} & \mathbb{Z}) \end{array}$$

The diagram shows a commutative square with dashed arrows s_1 and s_2 connecting the top and bottom rows. The central vertical arrow is labeled 1 .

A non-zero morphism in $\text{Adel}(\mathbf{A})$.

In particular, $(\mathbb{Z} \xrightarrow{2} \mathbb{Z} \xrightarrow{2} \mathbb{Z})$ is not isomorphic to the zero object.

Features of Adel(\mathbf{A})

Canonical embedding

$$\mathbf{A} \longrightarrow \text{Adel}(\mathbf{A})$$

$$A \mapsto (0 \rightarrow A \rightarrow 0)$$

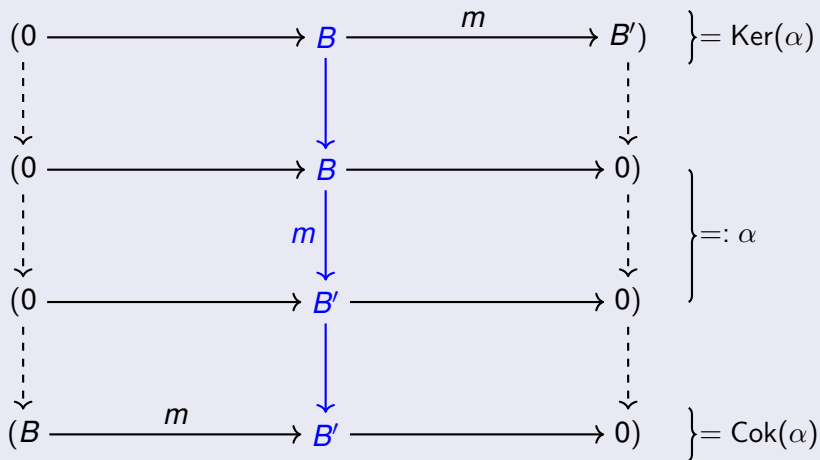
Duality

$$\text{Adel}(\mathbf{A}) \xrightarrow{\sim} \text{Adel}(\mathbf{A}^{\text{op}})^{\text{op}}$$

$$(A \xrightarrow{a} B \xrightarrow{b} C) \mapsto (C \xrightarrow{b^{\text{op}}} B \xrightarrow{a^{\text{op}}} A)$$

\rightsquigarrow if $\text{Adel}(\mathbf{A})$ has cokernels, then it also has kernels

How to compute (co)kernels



How to compute (co)kernels

$$\begin{array}{ccccc} (A \oplus A') & \xrightarrow{\begin{pmatrix} a & 0 \\ 0 & \text{id}_{A'} \end{pmatrix}} & B \oplus A' & \xrightarrow{\begin{pmatrix} b & m \\ 0 & a' \end{pmatrix}} & C \oplus B' \\ \downarrow & & \downarrow & & \downarrow \\ (A & \xrightarrow{a} & B & \xrightarrow{b} & C) \\ \downarrow & & \downarrow & & \downarrow \\ (A' & \xrightarrow{a'} & B' & \xrightarrow{b'} & C') \\ \downarrow & & \downarrow & & \downarrow \\ (A' \oplus B) & \xrightarrow{\begin{pmatrix} a' & 0 \\ m & b \end{pmatrix}} & B' \oplus C & \xrightarrow{\begin{pmatrix} b' & 0 \\ 0 & \text{id}_C \end{pmatrix}} & C' \oplus C \end{array} \left. \vphantom{\begin{array}{ccccc} (A \oplus A') & \xrightarrow{\begin{pmatrix} a & 0 \\ 0 & \text{id}_{A'} \end{pmatrix}} & B \oplus A' & \xrightarrow{\begin{pmatrix} b & m \\ 0 & a' \end{pmatrix}} & C \oplus B' \right\} = \text{Ker}(\alpha)$$

$$\left. \vphantom{\begin{array}{ccccc} (A & \xrightarrow{a} & B & \xrightarrow{b} & C) \\ (A' & \xrightarrow{a'} & B' & \xrightarrow{b'} & C') \end{array}} \right\} =: \alpha$$

$$\left. \vphantom{\begin{array}{ccccc} (A' \oplus B) & \xrightarrow{\begin{pmatrix} a' & 0 \\ m & b \end{pmatrix}} & B' \oplus C & \xrightarrow{\begin{pmatrix} b' & 0 \\ 0 & \text{id}_C \end{pmatrix}} & C' \oplus C \end{array}} \right\} = \text{Cok}(\alpha)$$

The diagram illustrates the computation of the kernel and cokernel of a linear map α . The map α is represented by the commutative diagram in the middle, where $(A \xrightarrow{a} B \xrightarrow{b} C)$ and $(A' \xrightarrow{a'} B' \xrightarrow{b'} C')$ are the components of α . The kernel of α is shown as the top row, and the cokernel is shown as the bottom row. The maps between the rows are represented by block matrices. The kernel is $\text{Ker}(\alpha)$ and the cokernel is $\text{Cok}(\alpha)$.

Adelman's theorem

Let \mathbf{A} be an additive category.

Theorem

$\mathbf{A} \hookrightarrow \text{Adel}(\mathbf{A})$ is the free abelian category of \mathbf{A} .

Its universal property is given as follows:

$$\begin{array}{ccc} \mathbf{A} & \begin{array}{c} \nearrow \\ \searrow \\ \xrightarrow{F} \end{array} & \begin{array}{c} \text{Adel}(\mathbf{A}) \\ \vdots \\ \mathbf{B} \end{array} \\ & & \end{array} \quad \begin{array}{c} (A \xrightarrow{a} B \xrightarrow{b} C) \\ \Downarrow \\ H(FA \xrightarrow{Fa} FB \xrightarrow{Fb} FC) \end{array}$$

Computability of Adelman categories

$\text{Adel}(\mathbf{A})$ provides a **computable** model of the free abelian category whenever we can solve 2-sided linear systems in \mathbf{A} .

1 Free abelian categories

2 A computational proof of the Snake lemma

Snake lemma

Suppose given the following commutative diagram with exact rows in an abelian category:

$$\begin{array}{ccccccc} A & \xrightarrow{a} & B & \longrightarrow & \text{Cok}(a) & \longrightarrow & 0 \\ & & \downarrow e & & \downarrow b & & \downarrow d \\ 0 & \longrightarrow & \text{Ker}(c) & \longrightarrow & C & \xrightarrow{c} & D \end{array}$$

Then we have an exact sequence

$$\text{Ker}(e) \longrightarrow \text{Ker}(b) \longrightarrow \text{Ker}(d) \xrightarrow{\partial} \text{Cok}(e) \longrightarrow \text{Cok}(b) \longrightarrow \text{Cok}(d)$$

- 1 Construct the universal abelian category \mathbf{U} which satisfies the premises of the Snake lemma.
- 2 Show that the conclusion holds for \mathbf{U} .

Snake lemma (rephrased)

Suppose given three consecutive morphisms in an abelian category

$$A \xrightarrow{a} B \xrightarrow{b} C \xrightarrow{c} D$$

such that

$$a \cdot b \cdot c = 0.$$

Then we have an exact sequence

$$\text{Ker}(e) \longrightarrow \text{Ker}(b) \longrightarrow \text{Ker}(d) \xrightarrow{\partial} \text{Cok}(e) \longrightarrow \text{Cok}(b) \longrightarrow \text{Cok}(d)$$

where d and e are defined by the following diagram:

$$\begin{array}{ccccccc} A & \xrightarrow{a} & B & \longrightarrow & \text{Cok}(a) & \longrightarrow & 0 \\ & & \downarrow e & & \downarrow b & & \downarrow d \\ 0 & \longrightarrow & \text{Ker}(c) & \longrightarrow & C & \xrightarrow{c} & D \end{array}$$

Construct the universal abelian category \mathbf{U} which satisfies the premises of the Snake lemma.

We set

$$\mathbf{U} := \text{Adel}(\mathbf{A})$$

where \mathbf{A} is the additive closure of the category freely generated by three consecutive morphisms $A \xrightarrow{a} B \xrightarrow{b} C \xrightarrow{c} D$ such that $a \cdot b \cdot c = 0$.

Example of a morphism in \mathbf{A}

$$A \oplus B \xrightarrow{\begin{pmatrix} \text{id}_A & a \cdot b \\ 0 & -3b \end{pmatrix}} A \oplus C$$

Remark

For every abelian category \mathbf{B} , an additive functor $\mathbf{A} \rightarrow \mathbf{B}$ is determined (up to nat. iso.) by three consecutive morphisms in \mathbf{B} whose composition yields 0.

Proof of the Snake lemma

Construct the universal abelian category \mathbf{U} which satisfies the premises of the Snake lemma.

We set

$$\mathbf{U} := \text{Adel}(\mathbf{A})$$




where \mathbf{A} is the additive closure of the category freely generated by three consecutive morphisms $A \xrightarrow{a} B \xrightarrow{b} C \xrightarrow{c} D$ such that $a \cdot b \cdot c = 0$.

Show that the conclusion holds for \mathbf{U} .

Software demo!

<https://github.com/sebastianpos/Adelman.jl>

References I

-  Murray Adelman, *Abelian categories over additive ones*, J. Pure Appl. Algebra **3** (1973), 103–117. MR 0318265
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